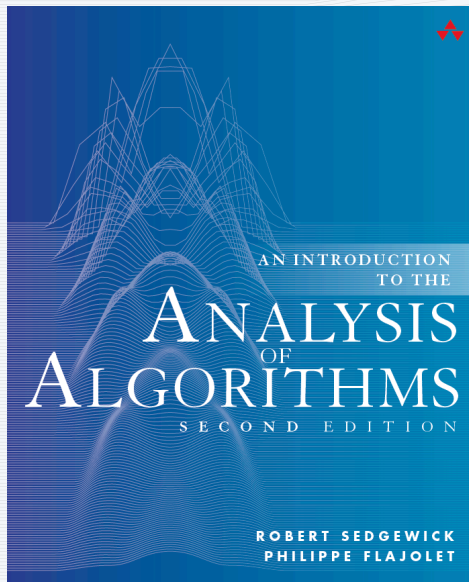


ANALYTIC COMBINATORICS

PART ONE



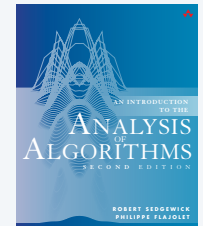
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8. Strings and Tries

Orientation

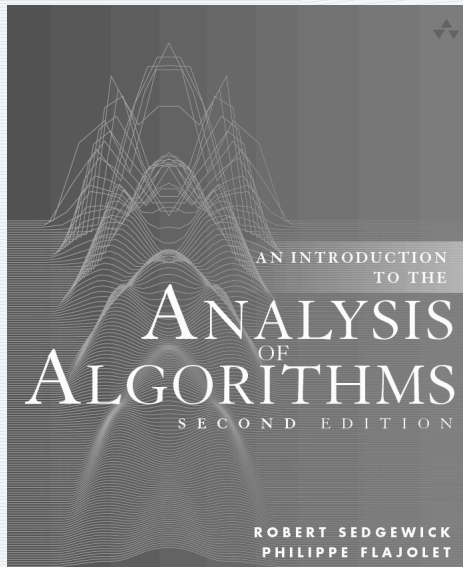
Second half of class

- Surveys fundamental combinatorial classes.
- Considers techniques from analytic combinatorics to study them .
- Includes applications to the analysis of algorithms.



<i>chapter</i>	<i>combinatorial classes</i>	<i>type of class</i>	<i>type of GF</i>
6	Trees	unlabeled	OGFs
7	Permutations	labeled	EGFs
8	Strings and Tries	unlabeled	OGFs
9	Words and Mappings	labeled	EGFs

Note: Many more examples in book than in lectures.



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8. Strings and Tries

- Bitstrings with restrictions
- Languages
- Tries
- Trie parameters

Bitstrings

```
23 10111110100101001100111000100111110110110100000111100001100111011101111101011000
9  110100101000111110100111100110100111011010111110000010110111001101000000111001110
29 1110111010110011101011100110100011001010111110011001000011001000101010010
6  1011100001101100011001110111001101101111011110011101011000011001100101000000110
13 1010110011101000110110111011001001011010010100110111100110000001111101000001111
1  1000001001100000110001100010000111100111001111000001100111110011011000100100111
24 1000101010111000111010110000011000001110101010001011000100110111110011110110010
18 0011101100101110010001100001001111010010011001100001100111010011010000101000111
42 00111111100110110111011011101010011011011100011111111010111010011000000100101110
5  1010100011110000101000001100100000110101001010001100110010101010111011011111110
2  11000000101111011011000101011010110010010000011101110010000001101010000000101000
70 11101111011011111011111111101001110100101111110111010011101001110100110010010
25 001111111001110101101111100010001000111000011101011110010101111100111010101111
0  00000010001111001110110101011100110000011110010010010101001100110011010011011110
24 1011110010100010011011110001100100011100100001010011010111011111010110010011100
7  01010010001011110110000110110101011010101111011001101101101000100110001111100111
23 01110110010011001110111000101010001101101001111111001101010111010001100110100001
3  00100011011010001100011111110011100110011110010110001100110011010001110111011101
```

Q. What is the *expected wait time for the first occurrence* of 000 in a random bitstring?

Q. What is the probability that an N -bit random bitstring *does not contain* 000?

Symbolic method for unlabelled objects (review)

Warmup: How many **binary strings** with N bits?

<i>Class</i>	B , the class of all binary strings
<i>Size</i>	$ b $, the number of bits in b
<i>OGF</i>	$B(z) = \sum_{b \in B} z^{ b } = \sum_{N \geq 0} B_N z^N$

Atoms

<i>type</i>	<i>class</i>	<i>size</i>	<i>GF</i>
0 bit	Z_0	1	z
1 bit	Z_1	1	z

Construction

$$B = \text{SEQ}(Z_0 + Z_1)$$

“a binary string is a sequence of 0 bits and 1 bits”

OGF equation

$$B(z) = \frac{1}{1 - 2z}$$

$$[z^N]B(z) = 2^N \quad \checkmark$$

Symbolic method for unlabelled objects (review)

Warmup: How many **binary strings** with N bits (alternate proof)?

<i>Class</i>	B , the class of all binary strings
<i>Size</i>	$ b $, the number of bits in b
<i>OGF</i>	$B(z) = \sum_{b \in B} z^{ b } = \sum_{N \geq 0} B_N z^N$

Atoms

<i>type</i>	<i>class</i>	<i>size</i>	<i>GF</i>
0 bit	Z_0	1	z
1 bit	Z_1	1	z

Construction

$$B = E + (Z_0 + Z_1) \times B$$

“a binary string is empty or a bit followed by a binary string”

OGF equation

$$B(z) = 1 + 2zB(z)$$

Solution

$$B(z) = \frac{1}{1 - 2z}$$

$$[z^N]B(z) = 2^N \quad \checkmark$$

Symbolic method for unlabelled objects (review)

Ex. How many N -bit binary strings have **no two consecutive 0s**?

<i>Class</i>	B_{00} , the class of binary strings with no 00
<i>OGF</i>	$B_{00}(z) = \sum_{b \in B_{00}} z^{ b }$

Construction $B_{00} = E + Z_0 + (Z_1 + Z_0 \times Z_1) \times B_{00}$

“a binary string with no 00 is either empty or 0 or it is 1 or 01 followed by a binary string with no 00”

OGF equation $B_{00}(z) = 1 + z + (z + z^2)B_{00}(z)$

Solution $B_{00}(z) = \frac{1 + z}{1 - z - z^2}$

Extract coefficients $[z^N]B_{00}(z) = F_N + F_{N+1} = F_{N+2}$ 1, 2, 5, 8, 13, ... ✓

$$= \frac{\phi^2}{\sqrt{5}} \phi^N \sim c_2 \beta_2^N \quad \text{with} \quad \begin{cases} \beta_2 \doteq 1.61803 \\ c_2 \doteq 1.17082 \end{cases}$$

Binary strings without long runs of 0s

Ex. How many N -bit binary strings have **no runs of P consecutive 0s**?

<i>Class</i>	B_P , the class of binary strings with no 0^P
<i>OGF</i>	$B_P(z) = \sum_{b \in B_P} z^{ b }$

Construction

$$B_P = Z_{<P}(E + Z_1 B_P)$$

OGF equation

$$B_P(z) = (1 + z + \dots + z^P)(1 + zB_P(z))$$

Solution

$$B_P(z) = \frac{1 - z^P}{1 - 2z + z^{P+1}}$$

Extract coefficients

$$[z^N]B_P(z) \sim c_k \beta_k^N \quad \text{where} \quad \begin{cases} \beta_k \text{ is the dominant root of } 1 - 2z + z^k \\ c_k = [\text{explicit formula available}] \end{cases}$$

See “Asymptotics” lecture

“a string with no 0^P is a string of 0s of length $<P$ followed by an empty string or a 1 followed by a string with no 0^P ”

Binary strings without long runs

Theorem. The number of binary strings of length N with no runs of P 0s is $\sim c_P \beta_P^N$ where c_P and β_P are easily-calculated constants.

```
sage: f2 = 1 - 2*x + x^3
sage: 1.0/f2.find_root(0, .99, x)
 $\beta_2$  1.61803398874989
sage: f3 = 1 - 2*x + x^4
sage: 1.0/f3.find_root(0, .99, x)
 $\beta_3$  1.83928675521416
sage: f4 = 1 - 2*x + x^5
sage: 1.0/f4.find_root(0, .99, x)
 $\beta_4$  1.92756197548293
sage: f5 = 1 - 2*x + x^6
sage: 1.0/f5.find_root(0, .99, x)
 $\beta_5$  1.96594823664510
sage: f6 = 1 - 2*x + x^7
sage: 1.0/f6.find_root(0, .99, x)
 $\beta_6$  1.98358284342432
```

Information on consecutive 0s in GFs for strings

$$S_P(z) = \sum_{s \in \mathcal{S}_P} z^{|s|} = \frac{1 - z^P}{1 - 2z + z^{P+1}} = \sum_{N \geq 0} \{\# \text{ of bitstrings of length } N \text{ with no } 0^P\} z^N$$

$$S_P(z/2) = \sum_{N \geq 0} (\{\# \text{ of bitstrings of length } N \text{ with no runs of } P \text{ 0s}\} / 2^N) z^N$$

$$S_P(1/2) = \sum_{N \geq 0} \{\# \text{ of bitstrings of length } N \text{ with no runs of } P \text{ 0s}\} / 2^N$$

$$= \sum_{N \geq 0} \Pr \{ \text{1st } N \text{ bits of a random bitstring have no runs of } P \text{ 0s} \}$$

$$= \sum_{N \geq 0} \Pr \{ \text{position of end of first } 0^P \text{ is } > N \} = \text{Expected position of end of first } 0^P$$

Theorem. Probability that an N -bit random bitstring has no 0^P : $[z^N] S_P(z/2) \sim c_P (\beta_P/2)^N$

Theorem. Expected wait time for the first 0^P in a random bitstring: $S_P(1/2) = 2^{P+1} - 2$

Consecutive 0s in random bitstrings

P	$S_P(z)$	approx. probability of no 0^P in N random bits			wait time
		N	10	100	
1	$\frac{1-z}{1-2z+z^2}$	$.5^N$	0.0010	$<10^{-30}$	2
2	$\frac{1-z^2}{1-2z+z^3}$	$1.1708 \times .80901^N$	0.1406	$<10^{-9}$	6
3	$\frac{1-z^3}{1-2z+z^4}$	$1.1375 \times .91864^N$	0.4869	0.0023	14
4	$\frac{1-z^4}{1-2z+z^5}$	$1.0917 \times .96328^N$	0.7510	0.0259	30
5	$\frac{1-z^5}{1-2z+z^6}$	$1.0575 \times .98297^N$	0.8906	0.1898	62
6	$\frac{1-z^6}{1-2z+z^7}$	$1.0350 \times .99174^N$	0.9526	0.4516	126

Validation of mathematical results

is **always** worthwhile when analyzing algorithms

```
public class TestOccP
{
    public static int find(int[] bits, int k)
        // See code at right.

    public static void main(String[] args)
    {
        int w = Integer.parseInt(args[0]);
        int maxP = Integer.parseInt(args[1]);
        int[] bits = new int[w];
        int[] sum = new int[maxP+1]; N/w trials.

        int T = 0;
        int cnt = 0;
        while (!StdIn.isEmpty())
        {
            T++;
            for (int j = 0; j < w; j++)
                bits[j] = BitIO.readbit();
            for (int P = 1; P <= maxP; P++)
                if (find(bits, P) == bits.length) sum[P]++;
        }

        for (int P = 1; P <= maxP; P++)
            StdOut.printf("%8.4f\n", 1.0*sum[P]/T);
        StdOut.println(T + " trials");
    }
}
```

- Read w-bits from StdIn
 - For each P, check for 0^P
- Print empirical probabilities.

```
public static int find(int[] bits, int P)
{
    int cnt = 0;
    for (int i = 0; i < bits.length; i++)
    {
        if (cnt == P) return i;
        if (bits[i] == 0) cnt++; else cnt = 0;
    }
    return bits.length;
}
```

```
% java TestOccP 100 6 < data/random1M.txt
0.0000    .0000
0.0000    .0000
0.0004    .0023
0.0267    .0259
0.1861    .1898
0.4502    .4516
10000 trials
```

← predicted by theory

✓

Wait time for specified patterns

```

9 23 10111110100101001100111000100111110110110100000111100001100111011101111101011000
4 9 110100101000111110100111100110100111011010111110000010110111001101000000111001110
12 29 1110111010110011101011100110100011000111001010111110011001000011001000101010010
8 5 1011100001101100011001110111001101101111011110011101011000011001100101000000110
6 13 1010110011101000110110111011001001011010010100110111100110000001111101000001111
4 1 10000010011000001100011000100001111001110011110000011001111110011011000100100111
2 24 1000101010111000111010110000011000001110101010001011000100110111110011110110010
0 18 0011101100101110010001100001001111010010011001100001100111010011010000101000111
0 42 0011111110011011011101101110101001101101110001111111010111010011000000100101110
6 5 1010100011110000101000001100100000110101001010001100110010101010111011011111110
6 2 11000000101111011011000101011010110010010000011101110010000001101010000000101000
30 70 111011110110111110111111111010011101001011111101110100111010011101001100010010010
0 25 0011111110011101011011111000100010001110000111010111100101011111001110101011111
4 0 00000010001111001110110101011100110000011110010010010101001100110011010011011110
6 24 1011110010100010011011110001100100011100100001010011010111011111010110010011100
4 7 01010010001011110110000110110101011010101111011001101101101000100110001111100111
7 23 01110110010011001110111000101010001101101001111111001101010111010001100110100001
0 3 00100011011010001100011111110011100110011110010110001100110011010001110111011101

```

Expected wait time for the first occurrence of 000: 17.9

Expected wait time for the first occurrence of 001: 6.0

Are these bitstrings random??

Autocorrelation

The probability that an N -bit random bitstring does not contain 0000 is $\sim 1.0917 \times .96328^N$

The expected wait time for the first occurrence of 0000 in a random bitstring is 30.

Q. Do the same results hold for 0001?

A. NO!

101111101001010011001110001001111101101101000001111100001

0001 occurs much
earlier than 0000

Observation. Consider first occurrence of 000.

- 0000 and 0001 equally likely, BUT
- mismatch for 0000 means 0001, so need to wait four more bits
- mismatch for 0001 means 0000, so *next* bit could give a match.

Q. What is the probability that an N -bit random bitstring does not contain 0001?

Q. What is the expected wait time for the first occurrence of 0001 in a random bitstring?

Constructions for strings without specified patterns

Cast of characters:

p — a pattern

S_p — binary strings that do not contain p

T_p — binary strings that *end in p*
and have no other occurrence of p

p 101001010

S_p 10111110101101001100110000011111

T_p 10111110101101001100110101001010

First construction

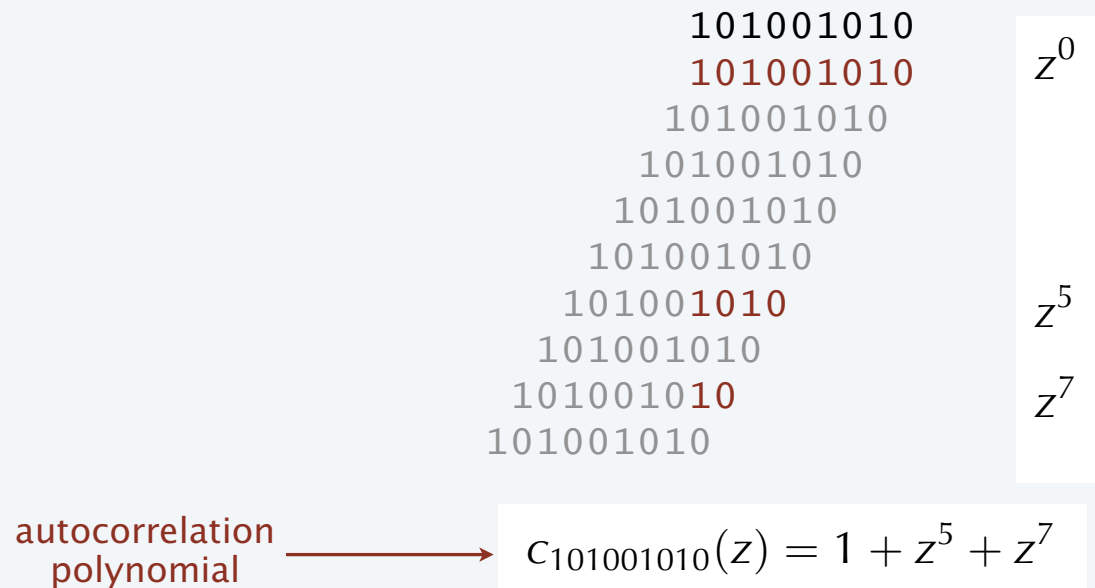
- S_p and T_p are disjoint
- the empty string is in S_p
- adding a bit to a string in S_p gives a string in S_p or T_p

$$S_p + T_p = E + S_p \times \{Z_0 + Z_1\}$$

Constructions for bitstrings without specified patterns

Every pattern has an autocorrelation polynomial

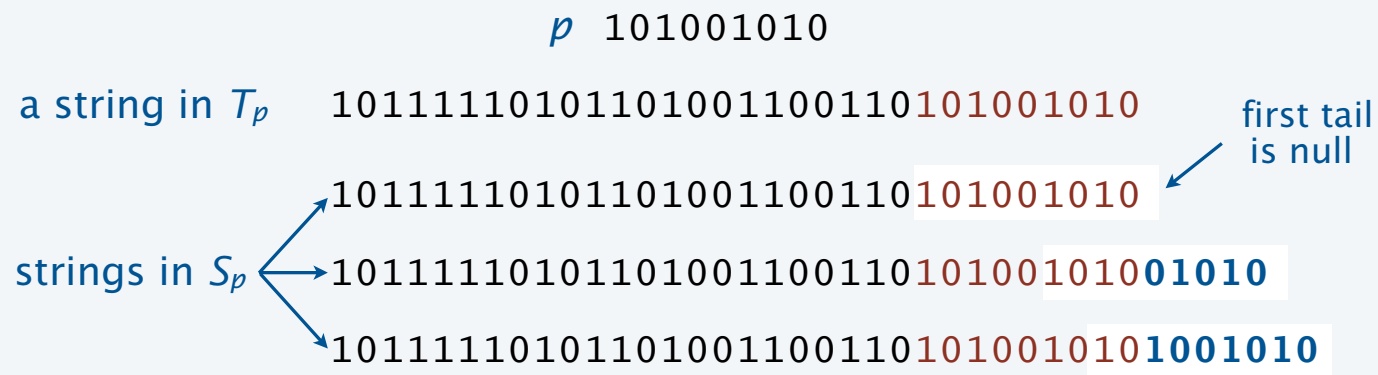
- slide the pattern to the left over itself.
- for each match of i trailing bits with the leading bits include a term $z^{|p| - i}$



Constructions for bitstrings without specified patterns

Second construction

- for each 1 bit in the autocorrelation of any string in T_p add a “tail”
- result is a string in S_p followed by the pattern



$$S_p \times \{p\} = T_p \times \sum_{c_i \neq 0} \{t_i\}$$

Bitstrings without specified patterns

How many N -bit strings **do not contain a specified pattern p** ?

<i>Classes</i>	S_p — the class of binary strings with no p
	T_p — the class of binary strings that end in p and have no other occurrence

<i>OGFs</i>	$S_p(z) = \sum_{s \in S_p} z^{ s }$
	$T_p(z) = \sum_{s \in T_p} z^{ s }$

Constructions

$$S_p + T_p = E + S_p \times \{Z_0 + Z_1\}$$

$$S_p \times \{p\} = T_p \times \sum_{c_i \neq 0} \{t_i\}$$

OGF equations

$$S_p(z) + T_p(z) = 1 + 2zS_p(z)$$

$$S_p(z)z^P = T_p(z)c_p(z)$$

Solution

$$S_p(z) = \frac{c_p(z)}{z^P + (1 - 2z)c_p(z)}$$

See “Asymptotics” lecture

Extract coefficients

$$[z^N]S_p(z) \sim c_p \beta_p^N \quad \text{where} \quad \begin{cases} \beta_p \text{ is the dominant root of } z^P + (1 - 2z)c_p(z) \\ c_p = [\text{explicit formula available}] \end{cases}$$

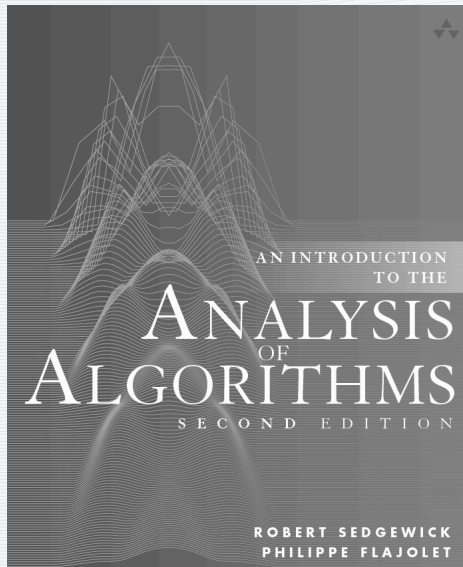
Autocorrelation for 4-bit patterns

p	auto-correlation	OGF	Probability that p does not occur in N random bits			wait time
			N	10	100	
0000 1111	1111	$\frac{1 - z^4}{1 - 2z + z^5}$	$.96328^N$	0.7510	0.0259	30
0001 0011 0111 1000 1100 1110	1000	$\frac{1}{1 - 2z + z^4}$	$.91964^N$	0.4327	0.0002	16
0010 0100 0110 1001 1011 1101	1001	$\frac{1 + z^3}{1 - 2z + z^3 - z^4}$	$.93338^N$	0.5019	0.0010	18
0101 1010	1010	$\frac{1 + z^2}{1 - 2z + z^2 - 2z^3 + z^4}$	$.94165^N$	0.5481	0.0024	20

Example. In 100 random bits,
 0000 is ~10 times more likely to be absent than 0101
 ~100 times more likely to be absent than 0001.

constants omitted
(close to 1)

off by < 10%
but indicative



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8. Strings and Tries

- Bitstrings with restrictions
- **Languages**
- Tries
- Trie parameters

Formal languages and the symbolic method

Definition. A **formal language** is a *set* of strings.

Q. How many strings of length N in a given language?

A. Use an OGF to enumerate them.

$$S(z) = \sum_{s \in \mathcal{S}} z^{|s|}$$

Remark. The symbolic method provides a systematic approach to this problem.

Issue. Ambiguity.

Regular expressions

Theorem. Let A and B be *unambiguous* REs with OGFs $A(z)$ and $B(z)$. If $A + B$, AB , and A^* are also unambiguous, then

$A(z) + B(z)$ enumerates $A + B$

$A(z)B(z)$ enumerates AB

$\frac{1}{1 - A(z)}$ enumerates A^*

OGF for an unambiguous RE is *rational* — can be written as the ratio of two polynomials.

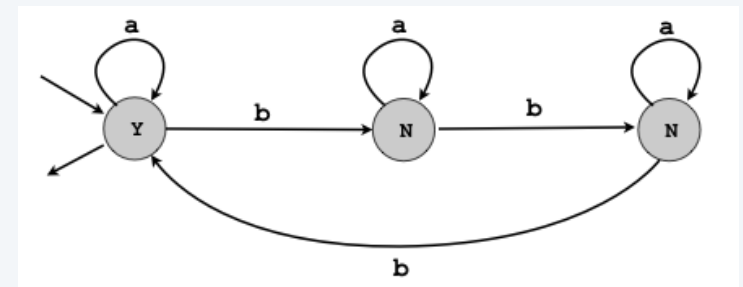
Proof.

Same as for symbolic method—different notation.

Corollary. OGFs that enumerate *regular languages* are rational.

Proof.

1. There exists an FSA for the language.
2. *Kleene's theorem* gives an unambiguous RE for the language defined by any FSA.



$a^* \mid (a^*ba^*ba^*ba^*)^*$

Regular expressions

Example 1. Binary strings with no 000

RE. $(1 + 01 + 001 + 001)^*(\epsilon + 0 + 00 + 00)$

OGF.
$$\begin{aligned} S_4(z) &= \frac{1 + z + z^2 + z^3}{1 - (z + z^2 + z^3 + z^4)} \\ &= \frac{\frac{1 - z^4}{1 - z}}{1 - z \frac{1 - z^4}{1 - z}} \\ &= \frac{1 - z^4}{1 - 2z + z^5} \quad \checkmark \end{aligned}$$

Expansion. $[z^N]S_4(z) \sim c_4 \beta_4^N$ with $\begin{cases} \beta_4 \doteq 1.92756 \\ c_4 \doteq 1.09166 \end{cases}$

Regular expressions

Example 2. Binary strings that represent multiples of 3

RE. $(1(01^*0)^*10^*)^*$

OGF.
$$D_3(z) = \frac{1}{1 - \frac{z^2}{1 - \frac{z^2}{1 - z}}} = \frac{1}{1 - \frac{z^2}{1 - z - z^2}}$$
$$= 1 - \frac{z^2}{(1 - 2z)(1 + z)}$$

Expansion. $[z^N]D_3(z) \sim \frac{2^{N-1}}{3} \quad \checkmark$

11
110
1001
1100
1111
10010
10101
11000
11011
11110
100001
100100
...

Context-free languages

Theorem. Let $\langle A \rangle$ and $\langle B \rangle$ be nonterminals in an *unambiguous* CFG with OGFs $A(z)$ and $B(z)$. If $\langle A \rangle \mid \langle B \rangle$ and $\langle A \rangle \langle B \rangle$ are also unambiguous, then

$A(z) + B(z)$ enumerates $\langle A \rangle \mid \langle B \rangle$

$A(z)B(z)$ enumerates $\langle A \rangle \langle B \rangle$

Proof.

Same as for symbolic method—different notation.

Corollary. OGFs that enumerate unambiguous CF languages are *algebraic*.

Proof.

"Gröbner basis" elimination—see text.

An *algebraic function* is a function that satisfies a polynomial equation whose coefficients are polynomials with rational coefficients

Context-free languages

The unlabelled constructions we have considered *are* CFGs, using different notation.

class	construction	CFG	OGF (algebraic)
Binary Trees	$T = E + T \times Z \times T$	$\langle T \rangle := \langle E \rangle$ $\langle T \rangle := \langle T \rangle \langle Z \rangle \langle T \rangle$	$T(z) = 1 + zT(z)^2$
Bitstrings	$B = E + (Z_0 + Z_1) \times B$	$\langle B \rangle := \langle E \rangle$ $\langle Y \rangle := \langle Z_0 \rangle \mid \langle Z_1 \rangle$ $\langle B \rangle := \langle Y \rangle \times \langle B \rangle$	$B(z) = 1 + 2zB(z)$
Bitstrings with no 00	$B_{00} = (E + Z_0)$ $\quad \times (E + Z_1 \times B_{00})$	$\langle Y_0 \rangle := \langle E \rangle \mid \langle Z_0 \rangle$ $\langle Y_1 \rangle := \langle Z_1 \rangle \times \langle B_{00} \rangle$ $\langle Y_2 \rangle := \langle E \rangle + \langle Y_1 \rangle$ $\langle B_{00} \rangle := \langle Y_0 \rangle \mid \langle Y_2 \rangle$	$B_{00}(z) = 1 + z$ $\quad + (z + z^2)B_{00}(z)$

Note 1. Not all CFGs correspond to combinatorial classes (ambiguity).

Note 2. Not all constructions are CFGs (many other operations have been defined).

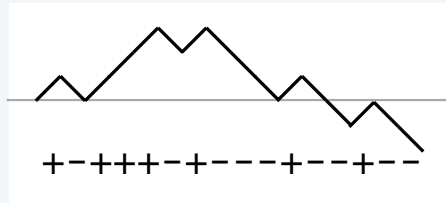
Walks

Definition. A **walk** is a sequence of + and – characters.

Sample applications:

- Parenthesis systems +-+++-+---
- Gambler's ruin problems
- Inversions in 2-ordered permutations (see text)

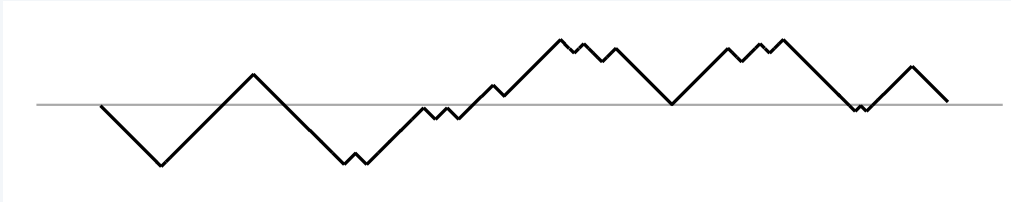
$()((()())())()$
 $+ - + + + - + - - - + - - + - -$



Q. How many different walks of length N ?

Q. How many different walks of length N where every prefix has more + than - ?

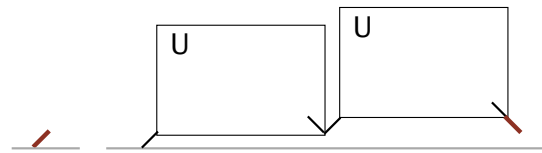
Unambiguous decomposition of walks



$\langle U \rangle$:

- start with +
- end at +1
- never hit 0

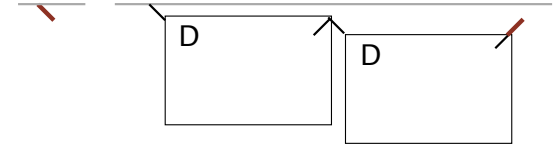
$$\langle U \rangle := \langle + \rangle \mid \langle U \rangle \langle U \rangle \langle - \rangle$$



<D>:

- start with $-$
- end at -1
- never hit 0

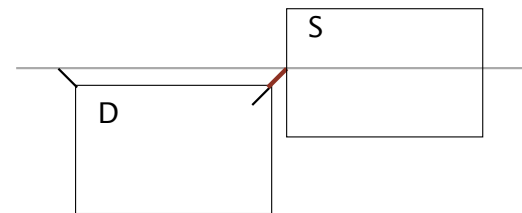
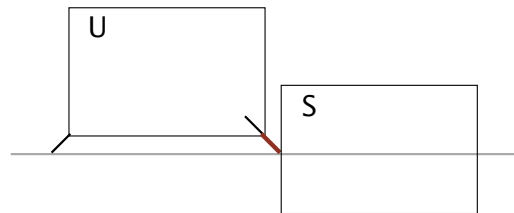
$$\langle D \rangle := \langle - \rangle \mid \langle D \rangle \langle D \rangle \langle + \rangle$$



<S>:

- begin at 0
- end at 0

$$\langle S \rangle := \langle U \rangle \langle - \rangle \langle S \rangle \mid \langle D \rangle \langle + \rangle \langle S \rangle$$



Context-free languages

Example. Walks of length $2N$ that start at and return to 0

CFL.

$$\begin{aligned}\langle S \rangle &:= \langle U \rangle \langle - \rangle \langle S \rangle \mid \langle D \rangle \langle + \rangle \langle S \rangle \mid \varepsilon \\ \langle U \rangle &:= \langle U \rangle \langle U \rangle \langle - \rangle \mid \langle + \rangle \\ \langle D \rangle &:= \langle D \rangle \langle D \rangle \langle + \rangle \mid \langle - \rangle\end{aligned}$$

OGFs.

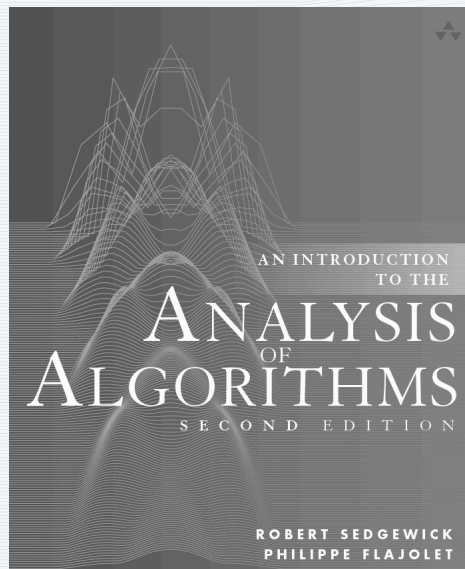
$$\begin{aligned}S(z) &= zU(z)S(z) + zD(z)S(z) + 1 \\ U(z) &= z + zU^2(z) \\ D(z) &= z + zD^2(z)\end{aligned}$$

Solve simultaneous equations.

$$\begin{aligned}U(z) = D(z) &= \frac{1}{2z} \left(1 - \sqrt{1 - 4z^2} \right) \\ S(z) &= \frac{1}{1 - 2zU(z)} = \frac{1}{\sqrt{1 - 4z^2}}\end{aligned}$$

Expand.

$$[z^{2N}]S(z) = \binom{2N}{N} \leftarrow \text{Elementary example, but extends to similar, more difficult problems}$$



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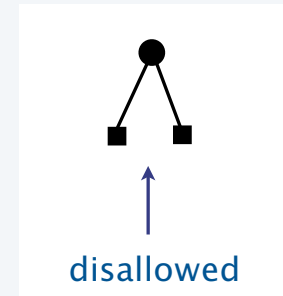
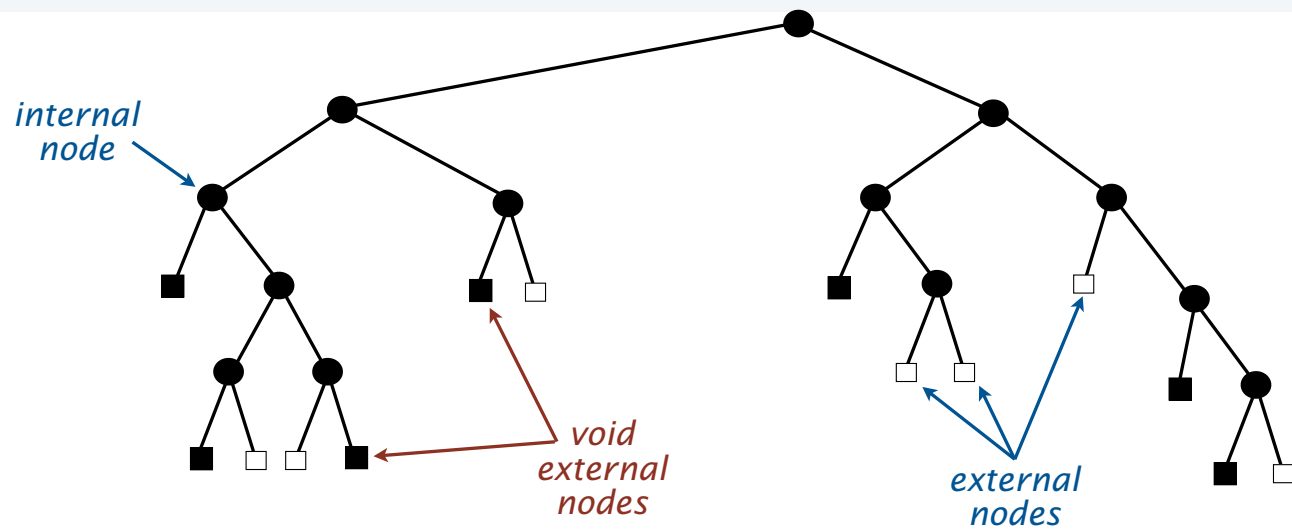
8. Strings and Tries

- Bitstrings with restrictions
- Languages
- **Tries**
- Trie parameters

Tries

Definition. A **trie** is a binary tree with the following properties:

- External nodes may be **void** (■)
- Siblings of void nodes are *not* void (● or □).

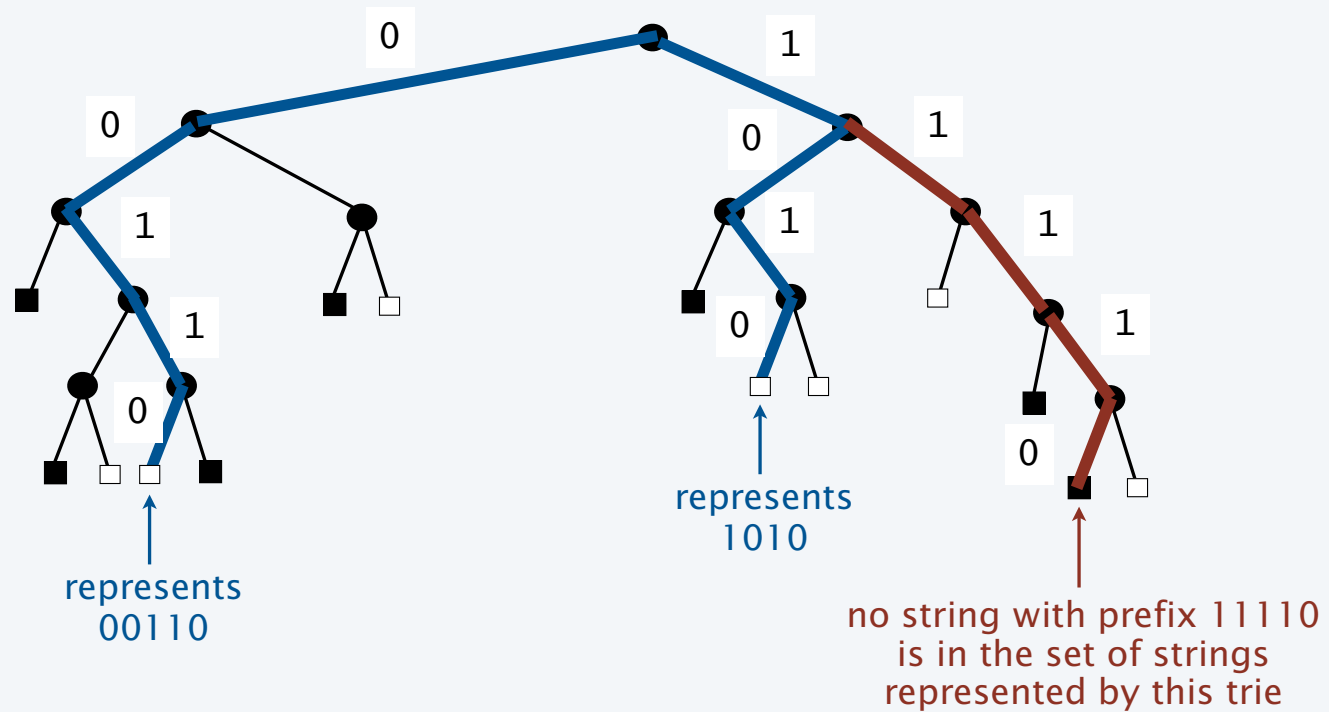


Ex. Give a recursive definition.

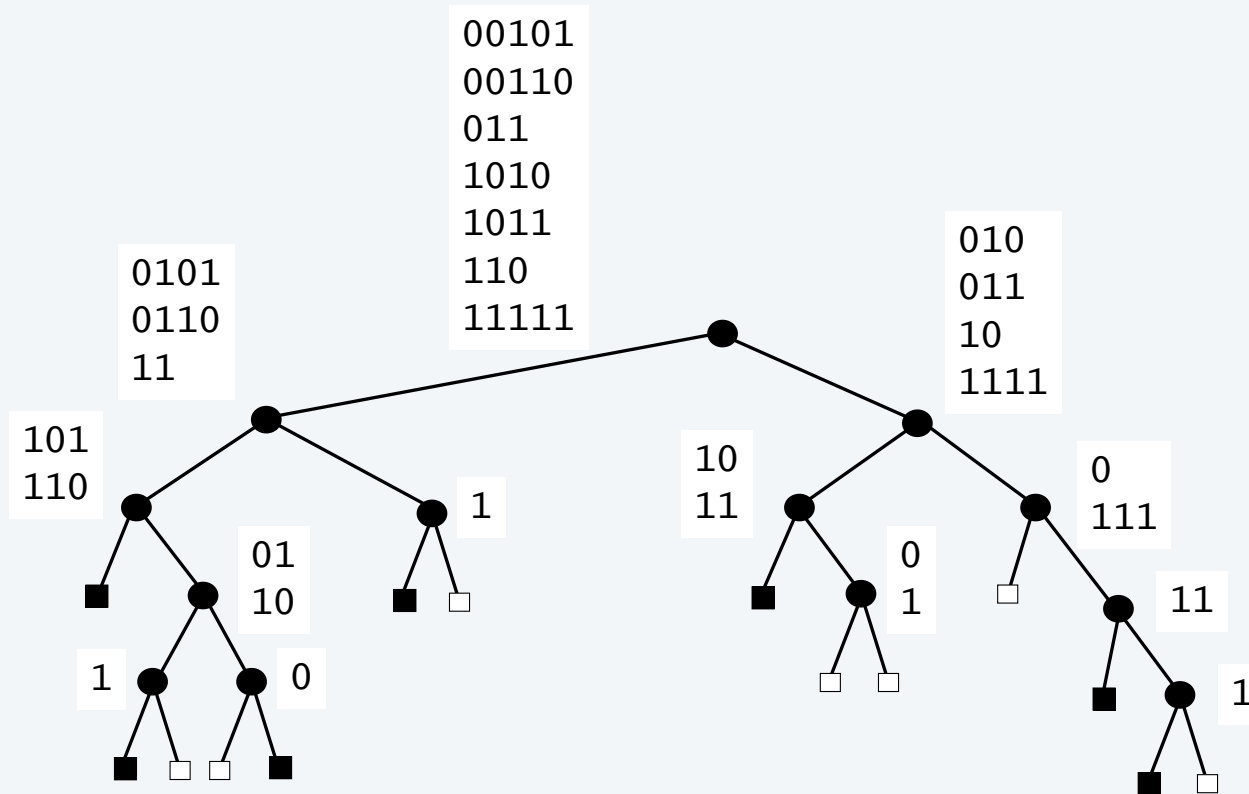
Tries and sets of bitstrings

Each trie corresponds to a set of bitstrings.

- Each nonvoid external node represents one bitstring.
- Path from the root to a node defines the bitstring



Tries and sets of bitstrings



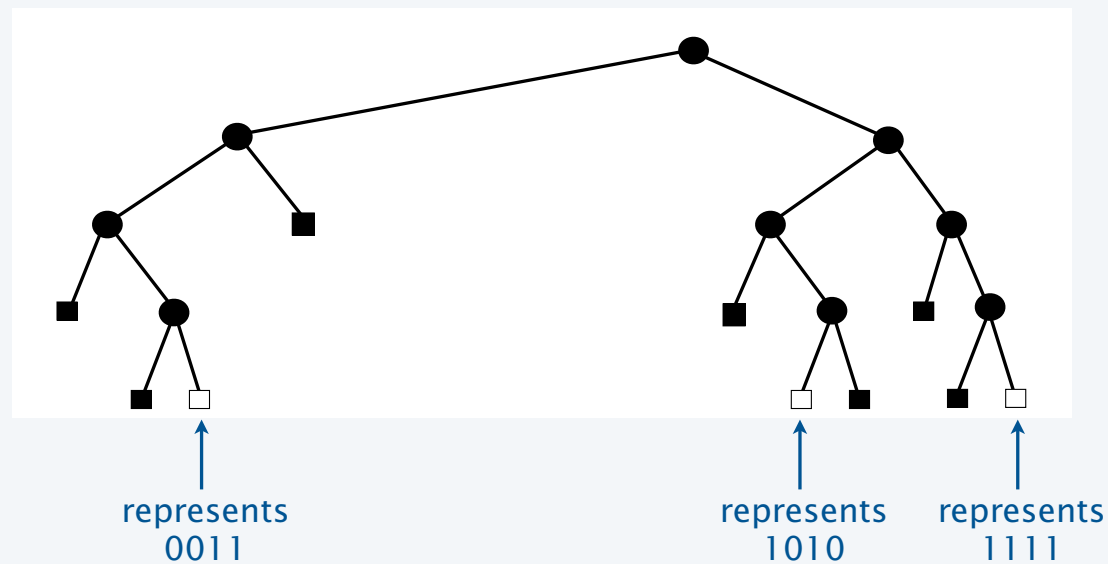
Note: Works only for *prefix-free* sets of bitstrings (or use void/nonvoid *internal* nodes).

no member is a prefix of another

Tries and sets of bitstrings (fixed length)

If all the bitstrings in the set are the same length, it is prefix-free.

0011
1010
1111



Trie applications

Searching and sorting

- MSD radix sort
- Symbol tables with string keys
- Suffix arrays

Data compression

- Huffman and prefix-free codes
- LZW compression

Decision making

- Collision resolution
- Leader election

Application areas:

Network systems

Bioinformatics

Internet search

Commercial data processing

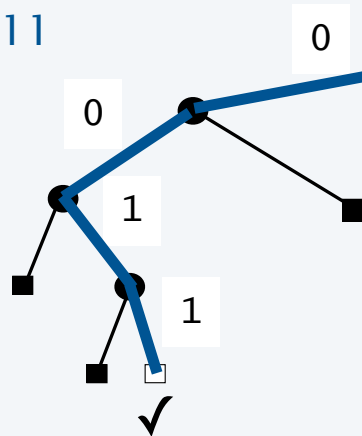


Trie application 1: Symbol tables

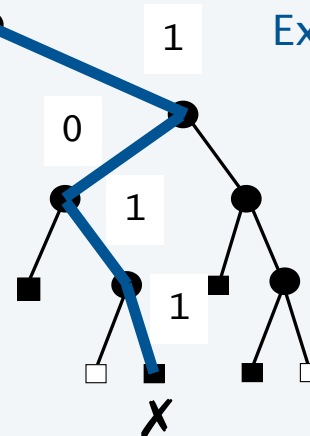
Search

- If at nonvoid external node and no bits left in bitstring, report success.
- If at void external node, report failure.
- If leading bit is 0, search in the left subtree (using remainder of string).
- If leading bit is 1, search in the right subtree (using remainder of string).

Ex: search for 0011



Ex: search for 10110



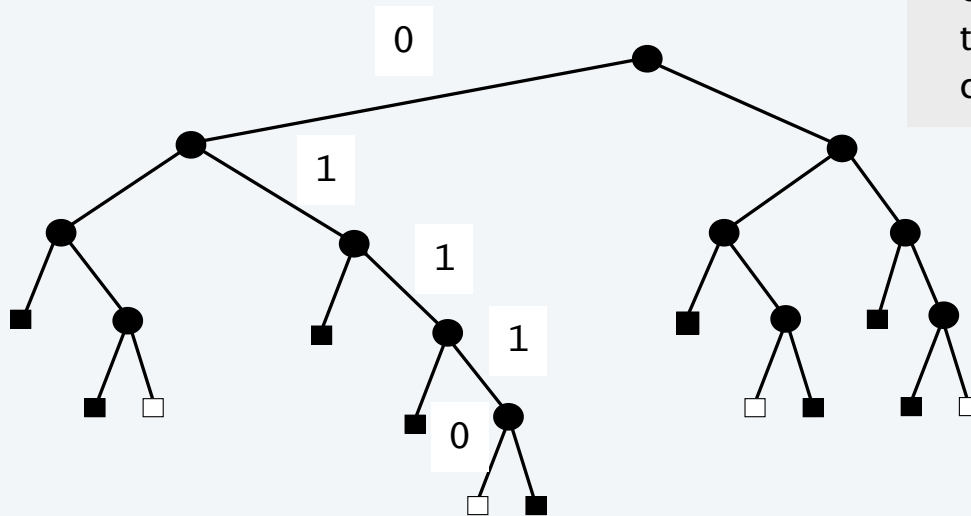
Q. Expected search time ?

Trie application 1: Symbol tables

Insert

- Search to void external node (prefix-free violation if nonvoid external node hit).
- Add internal nodes (each with one void external child) for each remaining bit.

Ex: insert 01110



variant:

convert the void external node
to a nonvoid external node that
contains a pointer to the "tail"

Q. How many void nodes ?

Trie application 2: Substring search index

Problem: Build an index that supports fast *substring search* in a given string S .

Ex. $S \rightarrow$

0	1	2	3	4	5	6	7	8	9
A	C	C	T	A	G	G	C	C	T

Q. Is ACCTA in S ?

A. Yes, starting at 0.

Q. Is CCT in S ?

A. Yes, in multiple places.

Q. Is TGA in S ?

A. No.

Solution: Use a *suffix multiway trie*.

Application 1: Search in genomic data.



Application 2: Internet search.



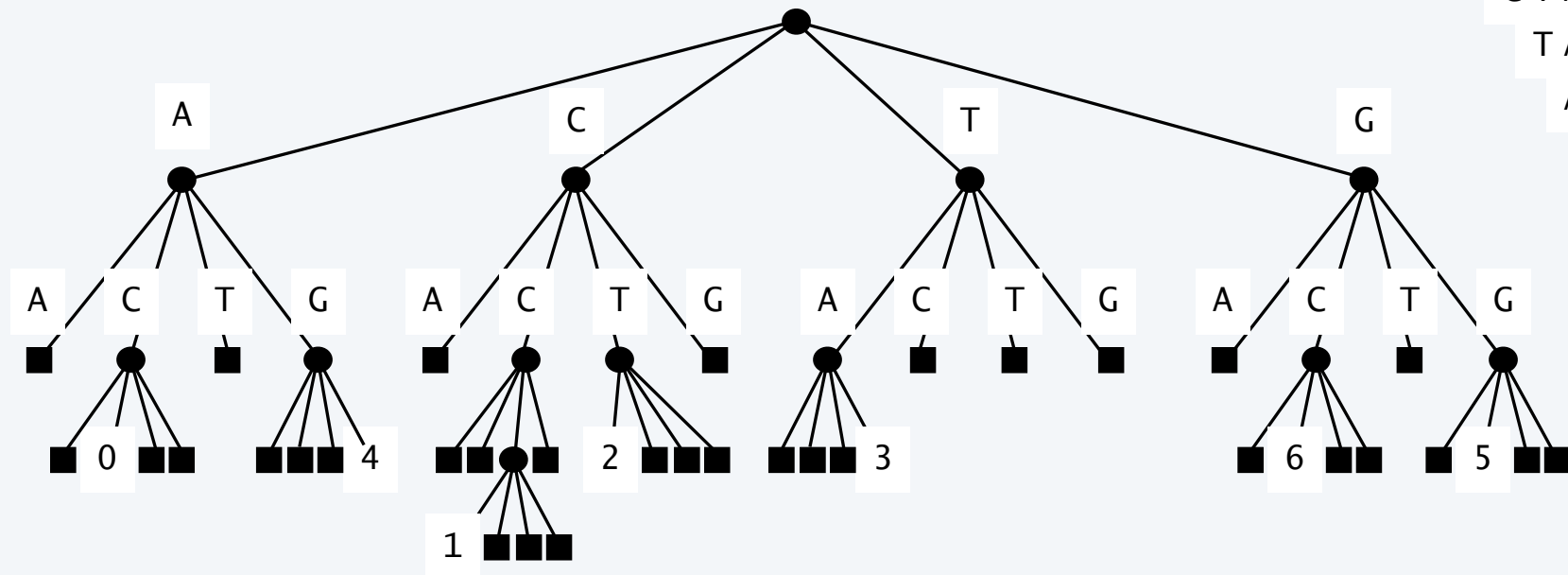
Trie application 2: Substring search index

To build the *suffix multiway trie* associated with a string S

- Insert the substrings starting at each position into an initially empty trie.
- Associate a string index with each nonvoid external node.

a prefix-free set

Property: Every internal node corresponds to a substring of S



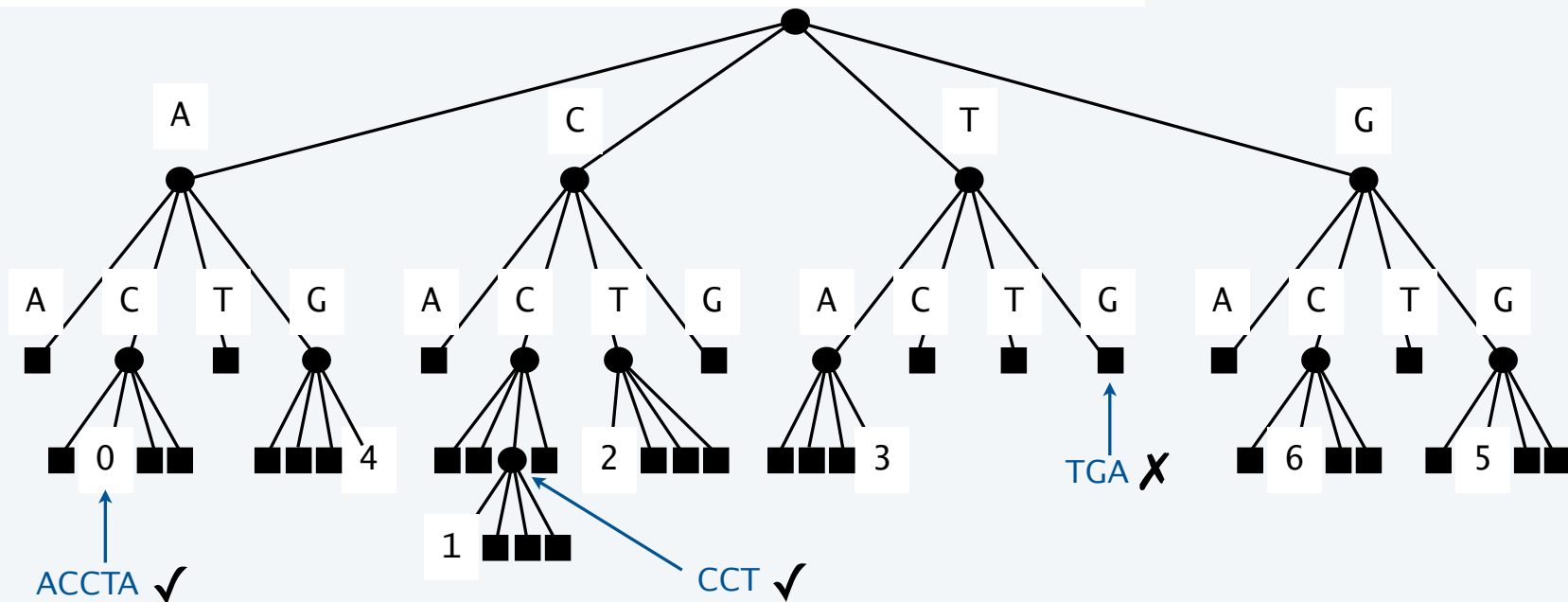
0	1	2	3	4	5	6	7	8	9
A	C	C	T	A	G	G	C	C	T
	C	C	T	A	G	G	C	C	T
		C	T	A	G	G	C	C	T
			T	A	G	G	C	C	T
				A	G	G	C	C	T
					G	G	C	C	T
						G	C	C	T
							C	C	T
								C	T
									T

Trie application 2: Substring index

To use a suffix tree to answer the query *Is X a substring of S?*

- Use the characters of *X* to traverse the trie.
- Continue in string when nonvoid node encountered.
- Report failure if void node encountered.
- Report success when end of *X* reached.

0	1	2	3	4	5	6	7	8	9
A	C	C	T	A	G	G	C	C	T

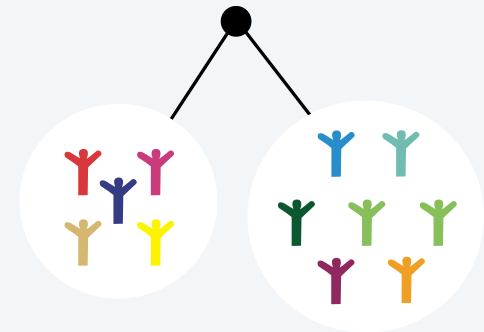
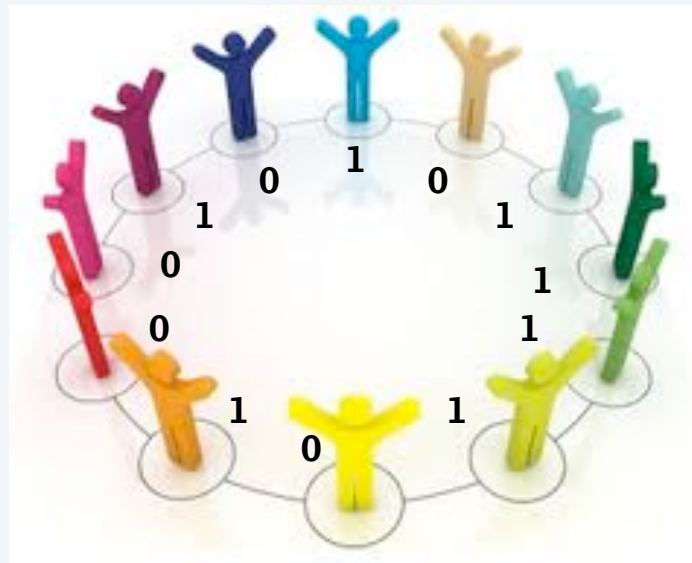


Trie application 3: Elect a leader



Problem: Elect a *leader* among a group of individuals.

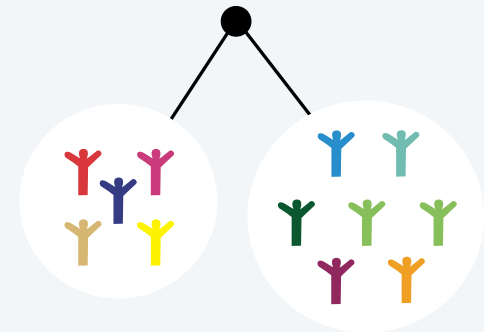
Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

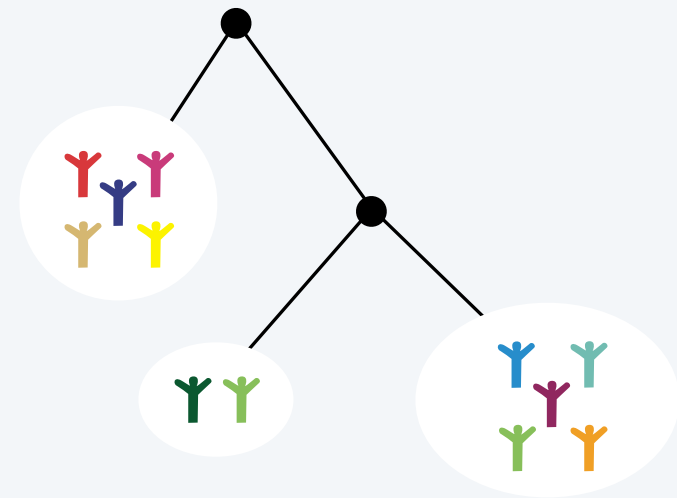
Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

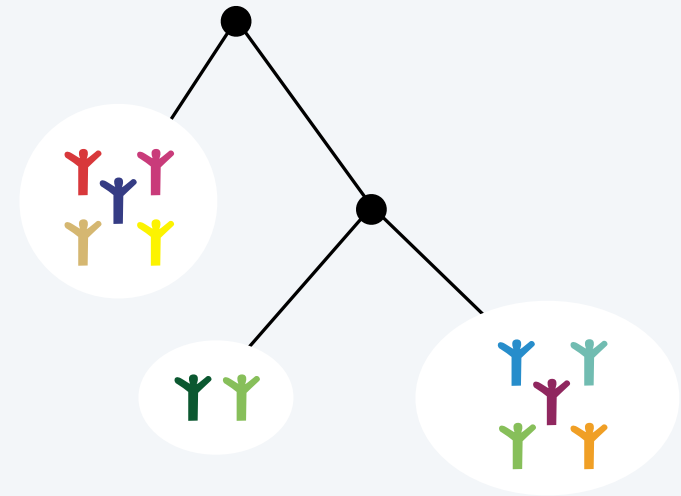
Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

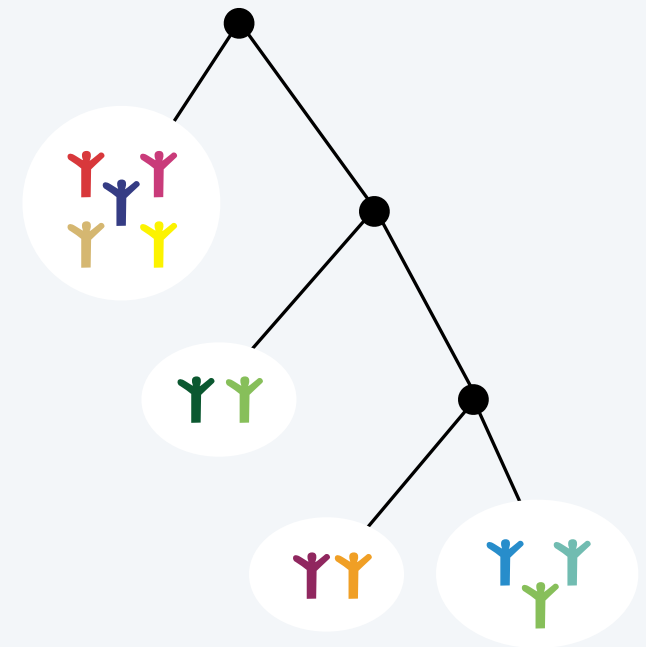
Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

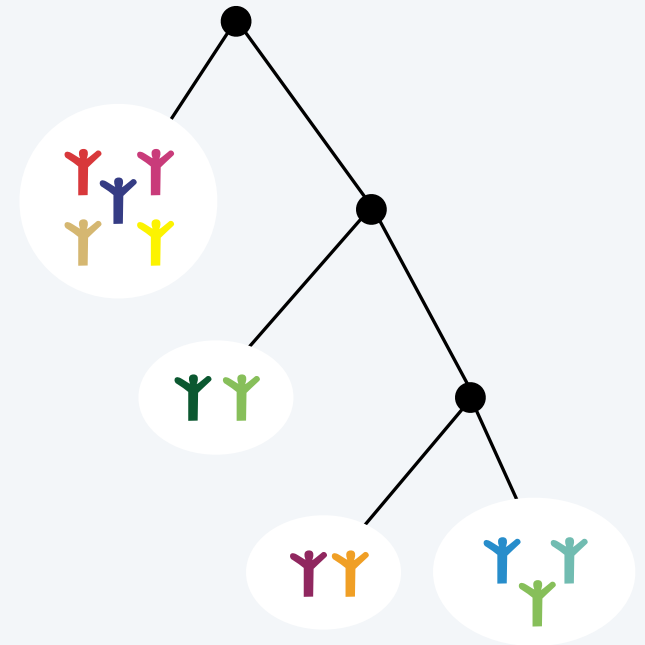
Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

Trie application 3: Elect a leader



Method.

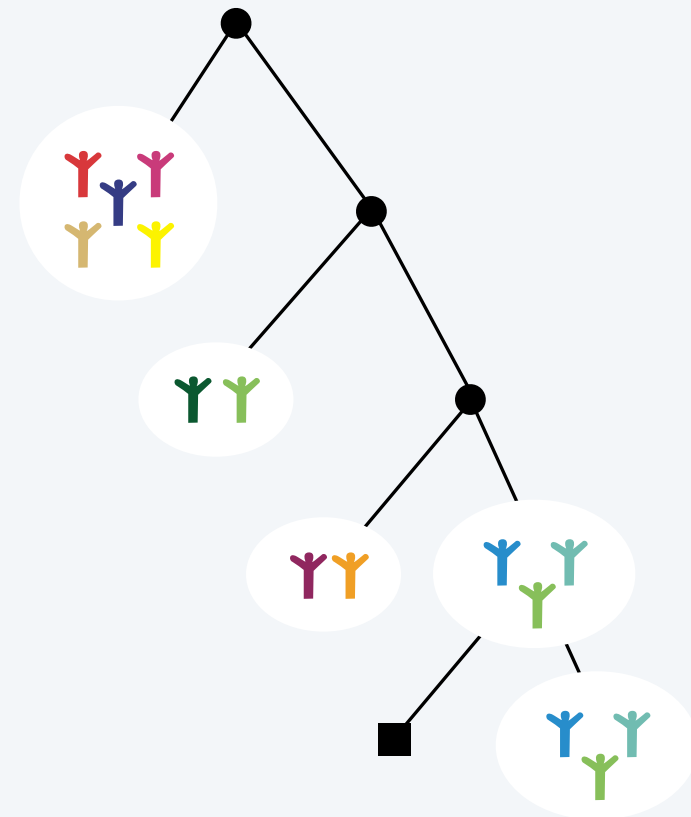
- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

Trie application 3: Elect a leader



Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

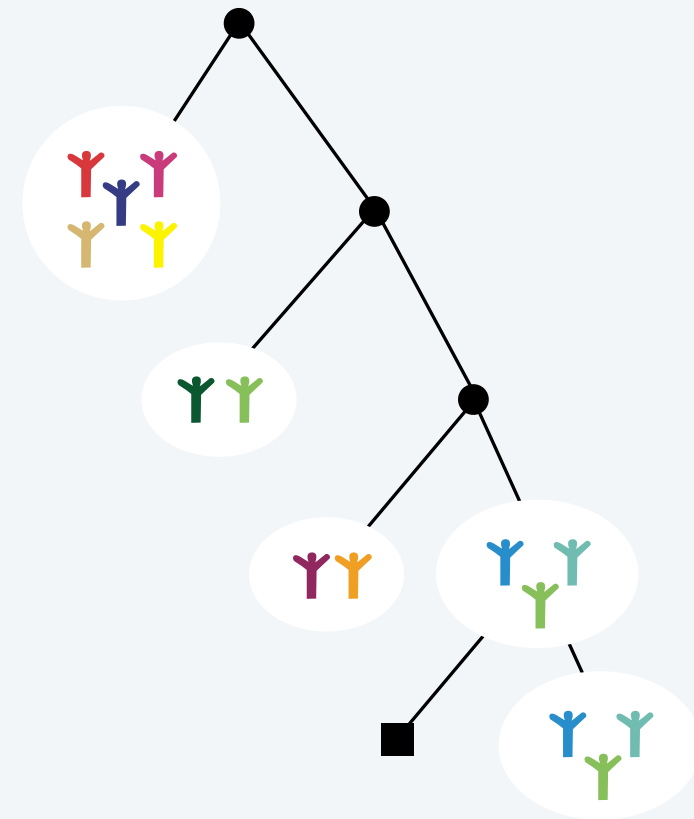


Trie application 3: Elect a leader

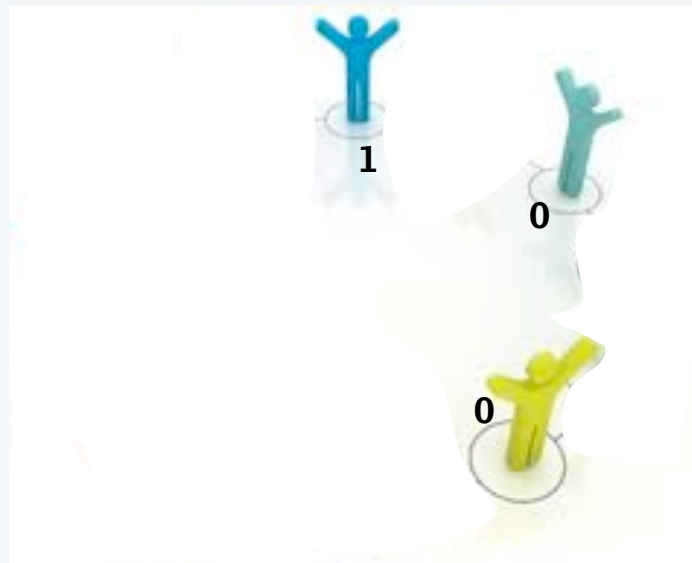


Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

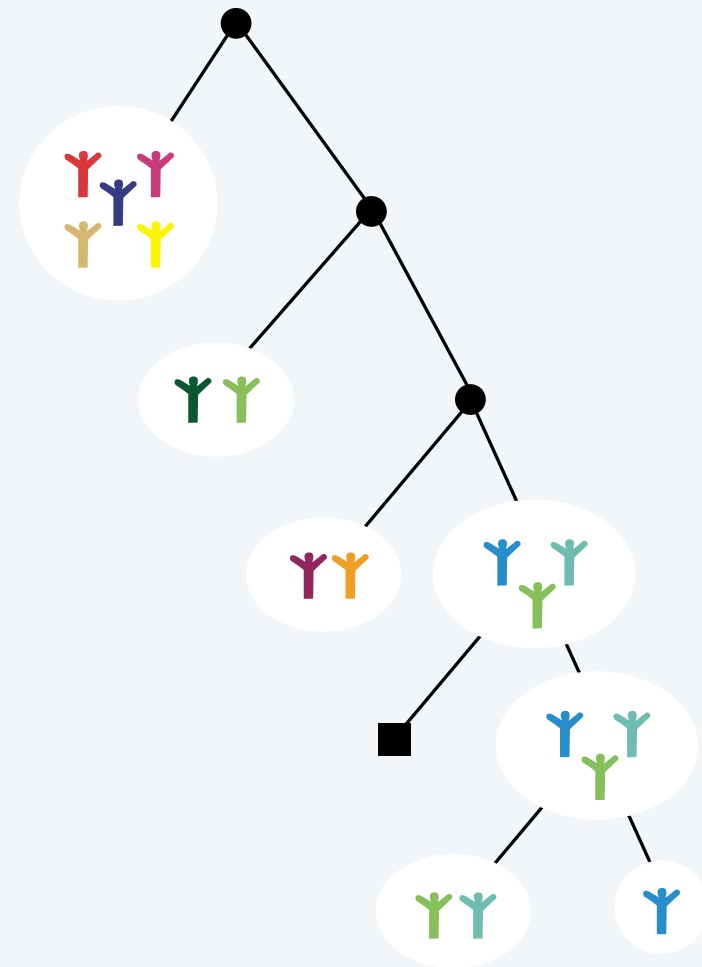


Trie application 3: Elect a leader

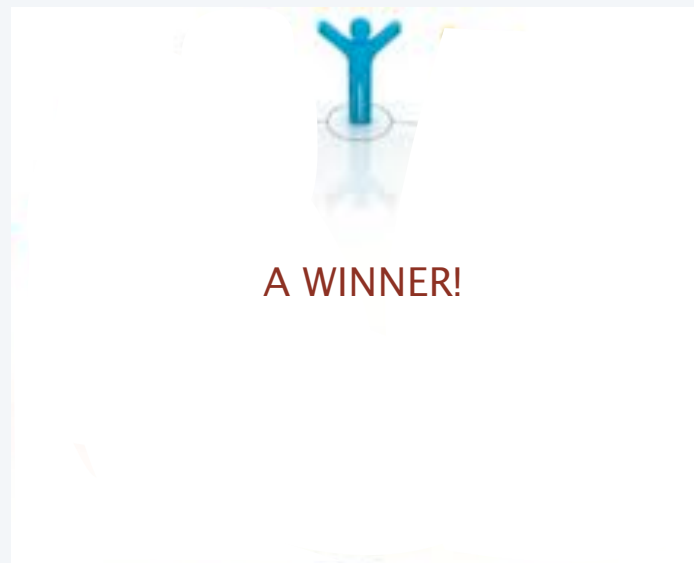


Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.

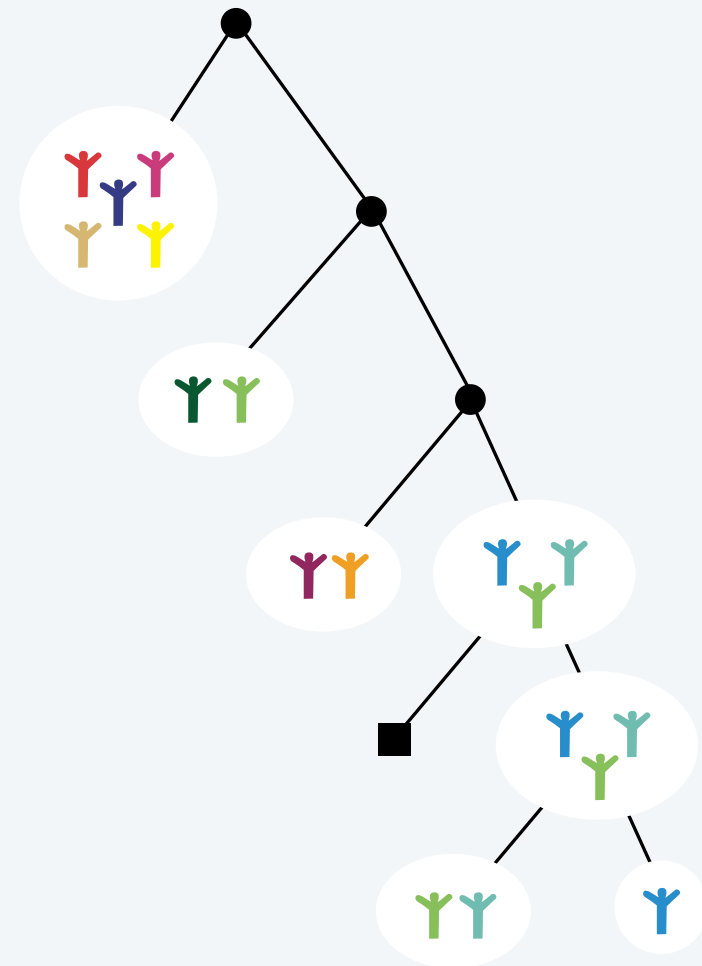


Trie application 3: Elect a leader



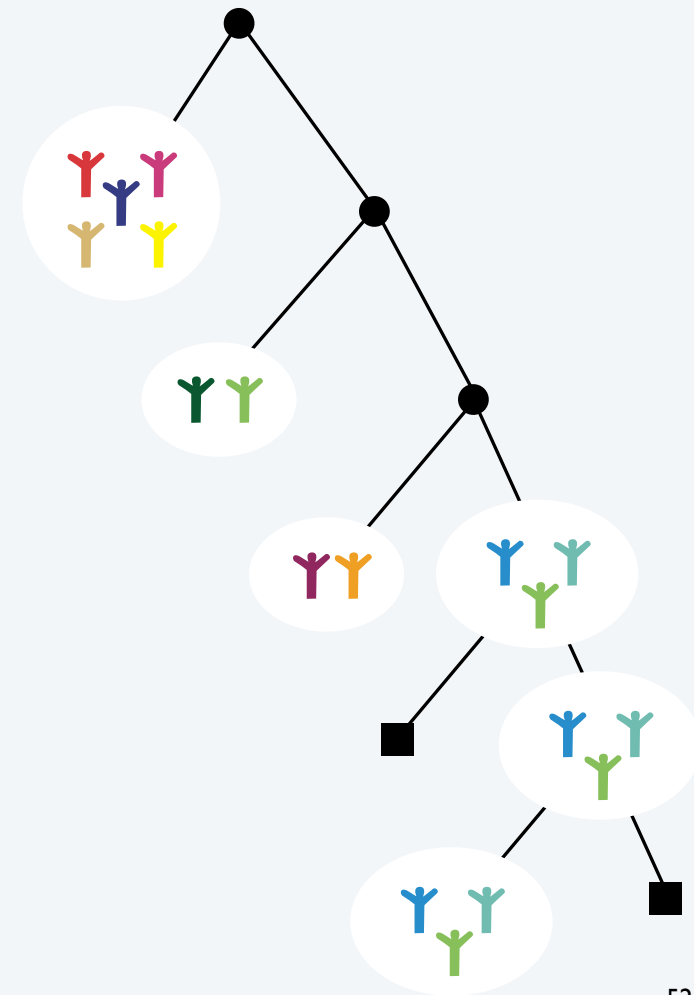
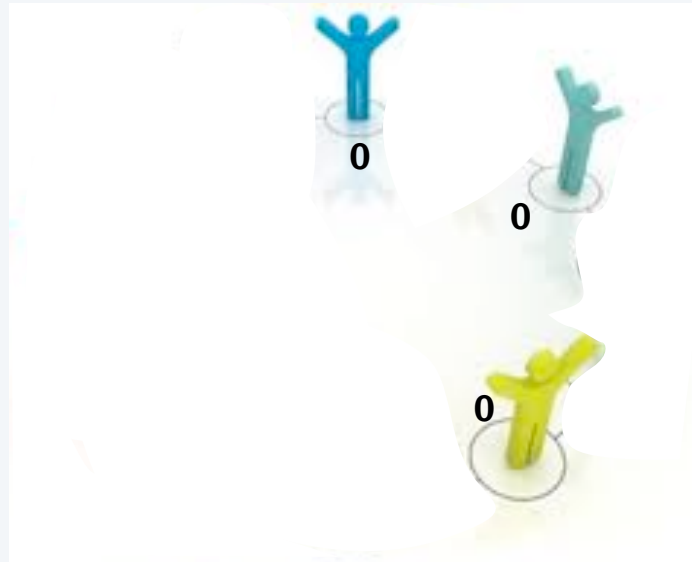
Method.

- Each person flips a 0-1 coin.
- 1 wins, 0 loses
- Winners continue to next round.



Trie application 3: Elect a leader

Procedure might fail!



Trie application 3: Elect a leader

a set of losers

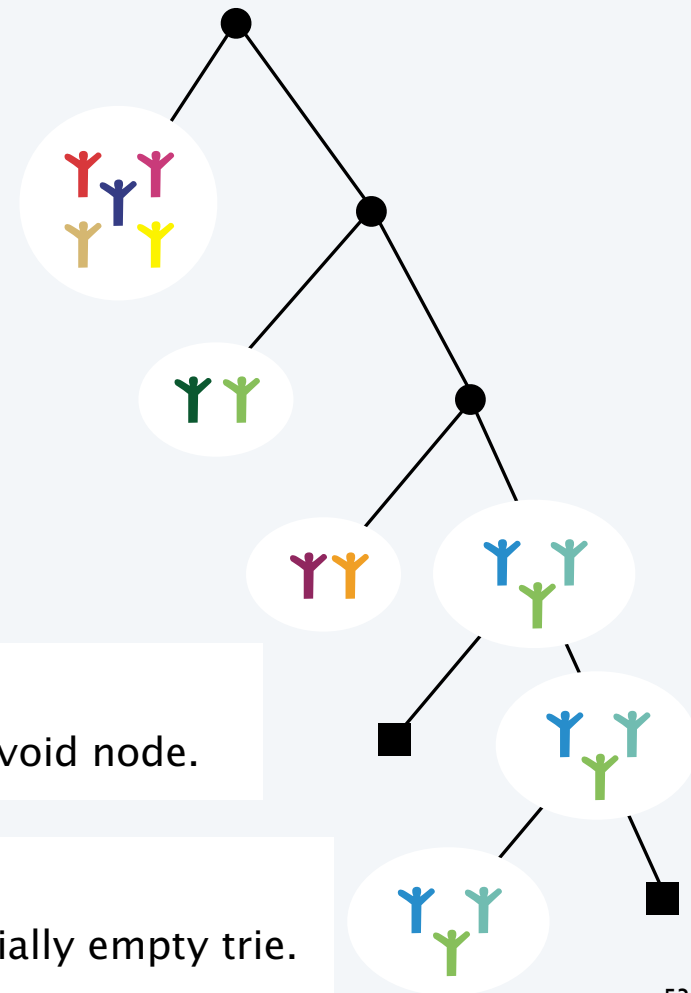
Procedure might fail!

Q. What is the chance of failure?

A. Probability that the rightmost path in a random trie ends in a void node.

Q. What is a random trie?

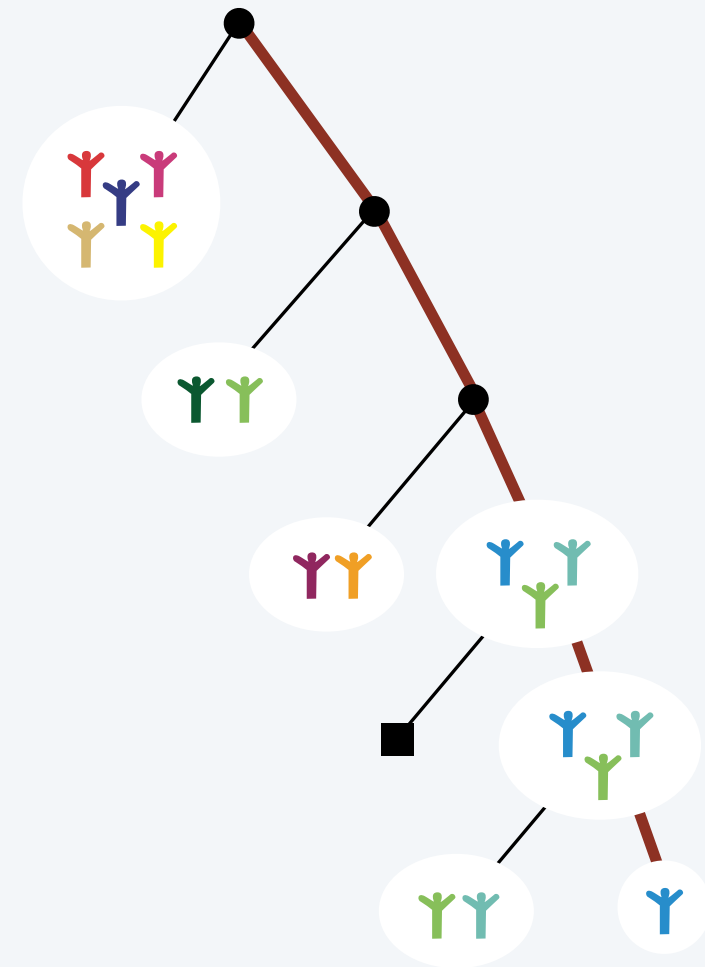
A. Built by inserting infinite-length random bitstrings into an initially empty trie.

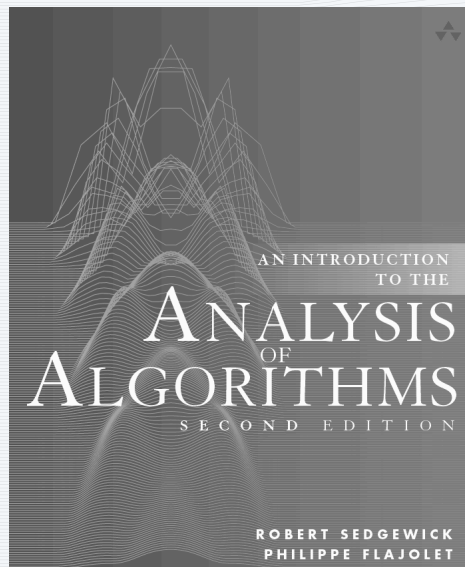


Trie application 3: Elect a leader



- Q. How many rounds in a distributed leader election?
A. Expected length of the rightmost path in a random trie.





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8. Strings and Tries

- Bitstrings with restrictions
- Languages
- Tries
- **Trie parameters**

`8d.Strings.TrieParms`

Analysis of trie parameters

is the basis of understanding performance in numerous large-scale applications.

Q. Space requirement?

A. Number of external nodes.

Q. "Extra" space ?

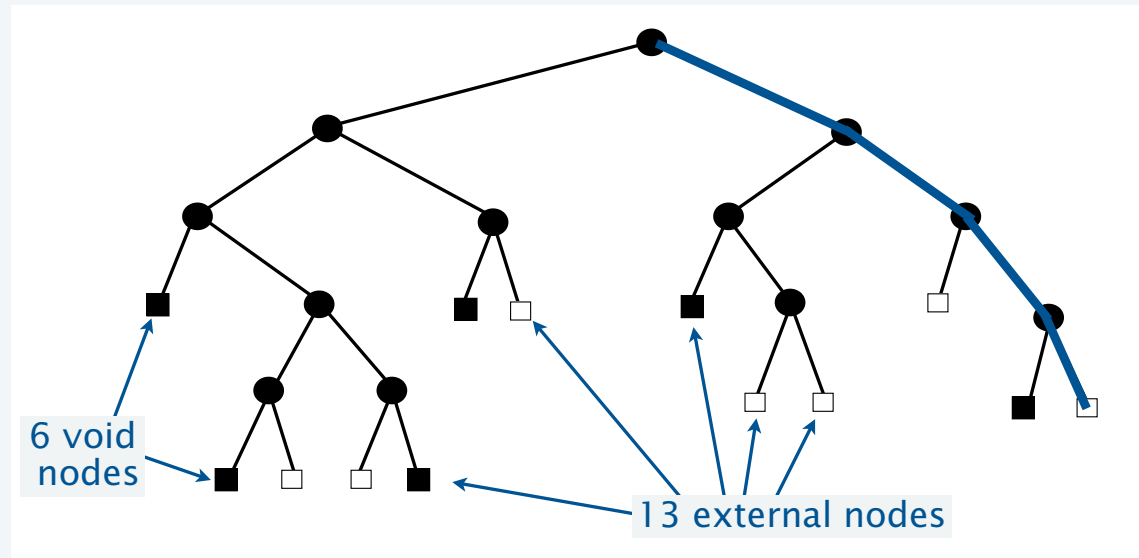
A. Number of void nodes.

Q. Expected search cost?

A. External path length.

Q. Rounds in leader election?

A. Length of rightmost path.

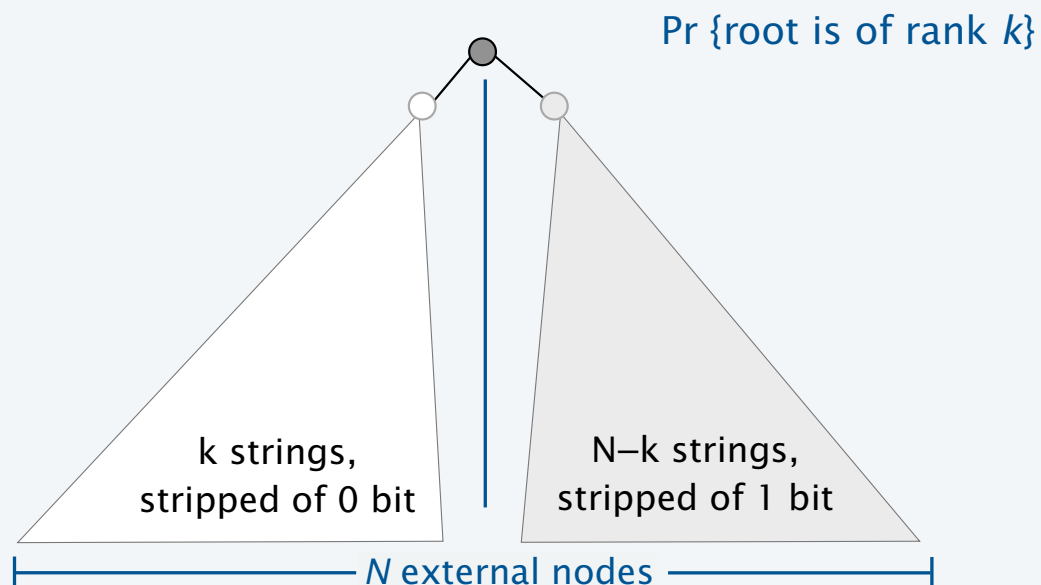


$$(3 + 5 + 5 + 5 + 5 + 3 + 3 + 3 + 4 + 4 + 3 + 4 + 4) / 13 \doteq 3.92$$

Usual model: Build trie from N *infinite* random bitstrings (nonvoid nodes represent tails)

Average external path length in a trie

Recurrence. [For comparison with BST and Catalan models.]



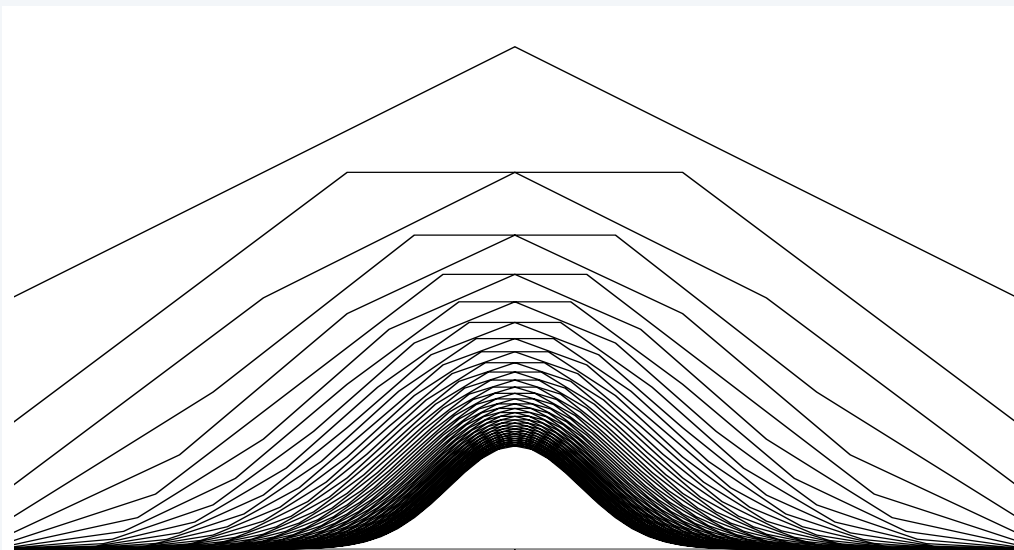
BST	$\frac{1}{N}$
Catalan	$\frac{\frac{1}{k} \binom{2k-2}{k} \frac{1}{N-k+1} \binom{2N-2k}{N-k}}{\frac{1}{N+1} \binom{2N}{N}}$
Trie	$\frac{1}{2^N} \binom{N}{k}$

$$C_N = N + \frac{1}{2^N} \sum_k \binom{N}{k} (C_k + C_{N-k}) \text{ for } N > 1 \text{ with } C_0 = C_1 = 0$$

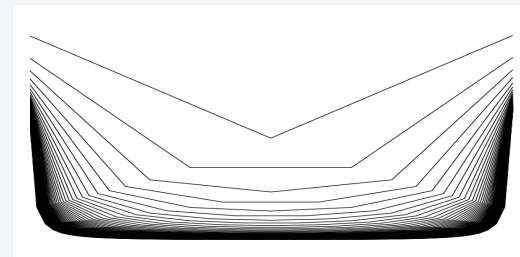
Caution: When $k = 0$ and $k = N$, C_N appears on right-hand side.

Probability that the root is of rank k in a random tree.

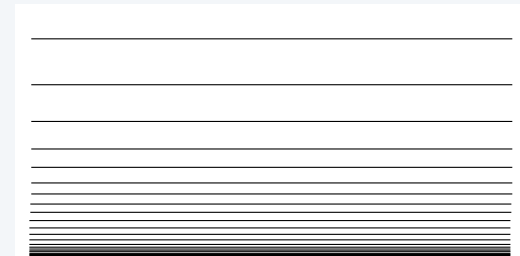
Trie built from random bitstrings



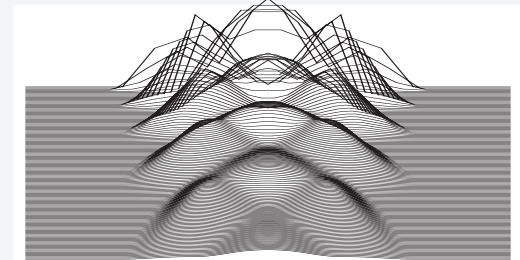
Random binary tree



BST built from random perm



AVL tree



Average external path length in a trie

Recurrence.

$$C_N = N + \frac{1}{2^N} \sum_k \binom{N}{k} (C_k + C_{N-k}) \text{ for } N > 1 \text{ with } C_0 = C_1 = 0$$

GF equation.

$$C(z) = ze^z - z + 2e^{z/2}C(z/2) \leftarrow \text{Also available directly through symbolic method}$$

EGF

$$C(z) = \sum_{N \geq 0} C_N \frac{z^N}{N!}$$

$$= ze^z - z + 2e^{z/2} \left(\frac{z}{2} e^{z/2} - \frac{z}{2} + 2e^{z/4}C(z/4) \right)$$

$$= z(e^z - 1) + z(e^z - e^{z/2}) + 4e^{3z/4}C(z/4)$$

$$= z(e^z - 1) + z(e^z - e^{z/2}) + z(e^z - e^{3z/4}) + 8e^{7z/8}C(z/8)$$

Iterate.

$$C(z) = z \sum_{j \geq 0} (e^z - e^{(1-2^{-j})z})$$

Expand.

$$C_N = N! [z^N] C(z) = N \sum_{j \geq 0} \left(1 - \left(1 - \frac{1}{2^j} \right)^{N-1} \right)$$

Approximate (exp-log)

$$C_N \sim N \sum_{j \geq 0} (1 - e^{-N/2^j}) \sim N \lg N \leftarrow \text{See next slide}$$

Average external path length in a trie

Goal: isolate periodic terms

$$\begin{aligned}\sum_{j \geq 0} (1 - e^{-N/2^j}) &= \sum_{0 \leq j < \lfloor \lg N \rfloor} (1 - e^{-N/2^j}) + \sum_{j \geq \lfloor \lg N \rfloor} (1 - e^{-N/2^j}) \\&= \lfloor \lg N \rfloor - \sum_{0 \leq j < \lfloor \lg N \rfloor} (e^{-N/2^j}) + \sum_{j \geq \lfloor \lg N \rfloor} (1 - e^{-N/2^j}) \\&= \lfloor \lg N \rfloor - \sum_{j < \lfloor \lg N \rfloor} (e^{-N/2^j}) + \sum_{j \geq \lfloor \lg N \rfloor} (1 - e^{-N/2^j}) + O(e^{-N}) \\&= \lfloor \lg N \rfloor - \sum_{j < 0} (e^{-N/2^{j+\lfloor \lg N \rfloor}}) + \sum_{j \geq 0} (1 - e^{-N/2^{j+\lfloor \lg N \rfloor}}) + O(e^{-N}) \\&= \lg N - \{\lg N\} - \sum_{j < 0} e^{-2^{\{\lg N\}-j}} + \sum_{j \geq 0} (1 - e^{-2^{\{\lg N\}-j}}) + O(e^{-N}) \quad \checkmark\end{aligned}$$

Average external path length in a trie

Q. $C_N = N + \frac{1}{2^N} \sum_k \binom{N}{k} (C_k + C_{N-k})$ for $N > 1$ with $C_0 = C_1 = 0$

A. $C_N/N = \lg N - \{ \lg N \} - \sum_{j < 0} e^{-2^{\{ \lg N \} - j}} + \sum_{j \geq 0} (1 - e^{-2^{\{ \lg N \} - j}}) + O(e^{-N})$

A

B

C

A



0.8

B



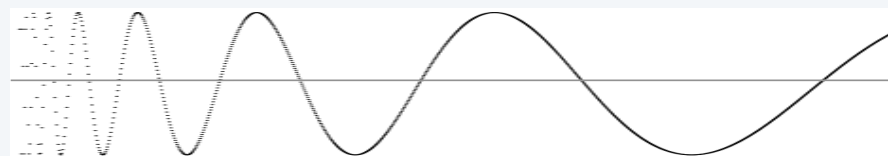
0.2

C



0.7

A + B + C

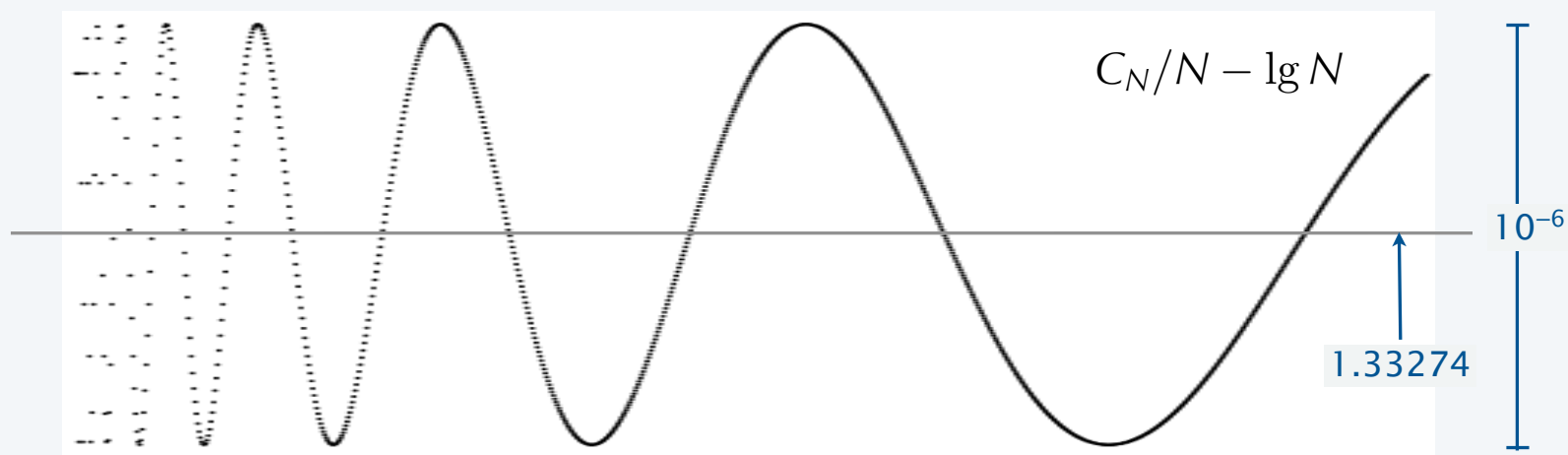


1.33274

10⁻⁶

Fluctuating term in trie (and other AofA) results

$$C_N = N + \frac{1}{2^N} \sum_k \binom{N}{k} (C_k + C_{N-k}) \text{ for } N > 1 \text{ with } C_0 = C_1 = 0$$

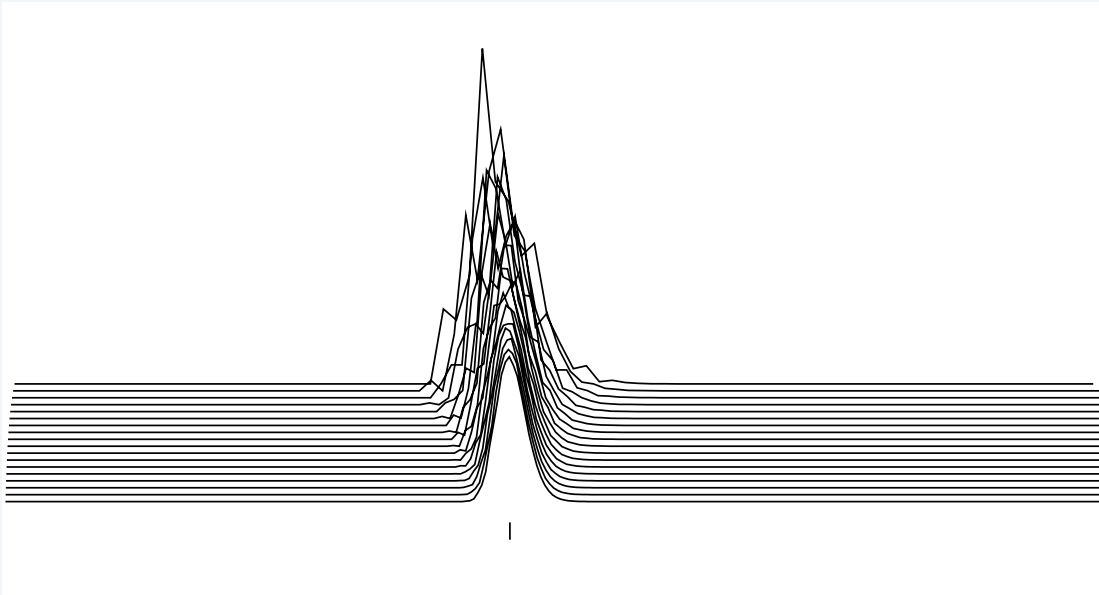


Q. Is there a reason that such a recurrence should imply such periodic behavior?

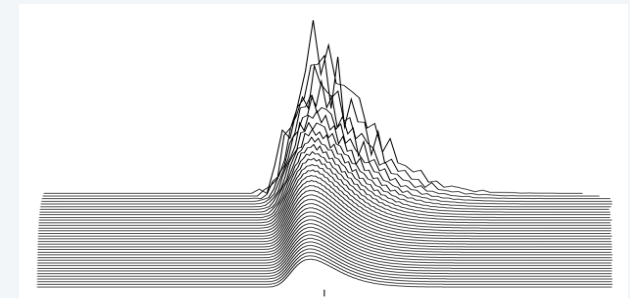
A. Yes. Stay tuned for the Mellin transform and related topics in Part II.

Average external path length distribution

Trie built from random bitstrings



BST built from random perm



Analysis of trie parameters

is the basis of understanding performance in numerous large-scale applications.

Q. Space requirement?

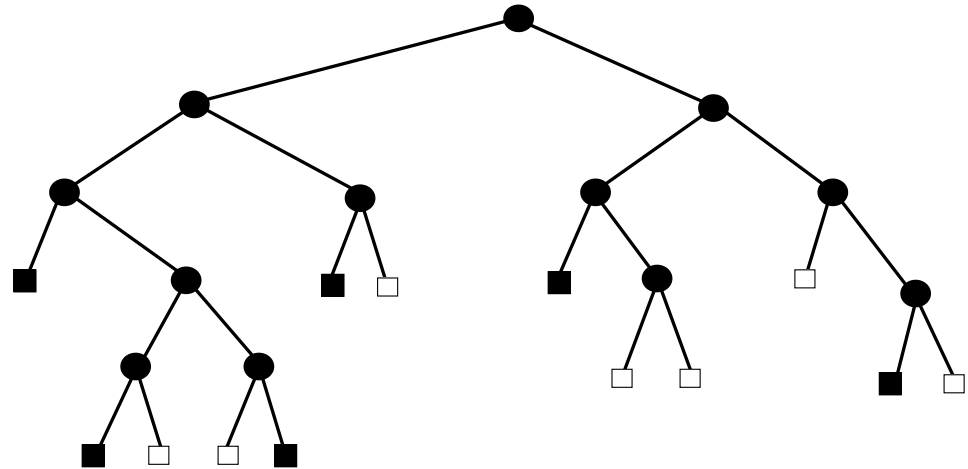
A. $\sim N/\ln 2 \doteq 1.44 N$.

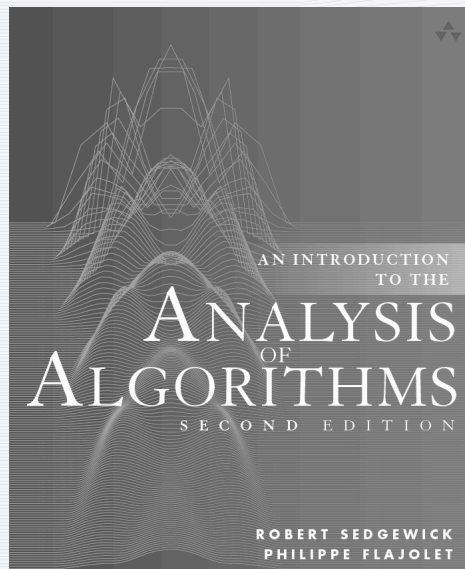
Q. Expected search cost?

A. About $N \lg N - 1.333 N$.

Q. Rounds in leader election?

A. [see exercise 8.57].





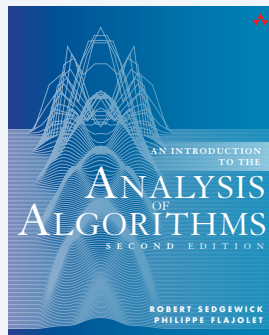
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8. Strings and Tries

- Bitstrings with restrictions
- Languages
- Tries
- Trie parameters
- **Exercises**

Exercise 8.3

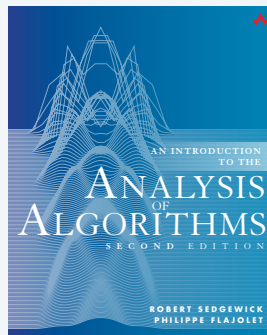
Good chance of a long run of 0s.



Exercise 8.3 How long a string of random bits should be taken to be 50% sure that there are at least 32 consecutive 0s?

Exercise 8.14

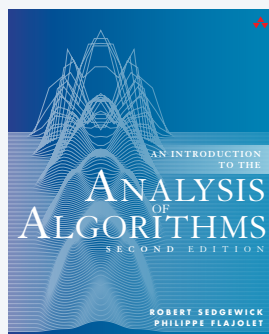
Monkey at a keyboard.



Exercise 8.14 Suppose that a monkey types randomly at a 32-key keyboard. What is the expected number of characters typed before the monkey hits upon the phrase
THE QUICK BROWN FOX JUMPED OVER THE LAZY DOG?

Exercise 8.57

Leader-election success probability.



Exercise 8.57 Solve the recurrence for p_N given in the proof of Theorem 8.9, to within the oscillating term.

$$p_N = \frac{1}{2^N} \sum_k \binom{N}{k} p_k \quad \text{for } N > 1 \text{ with } p_0 = 0 \text{ and } p_1 = 1$$

Assignments for next lecture

1. Read pages 415-472 in text.

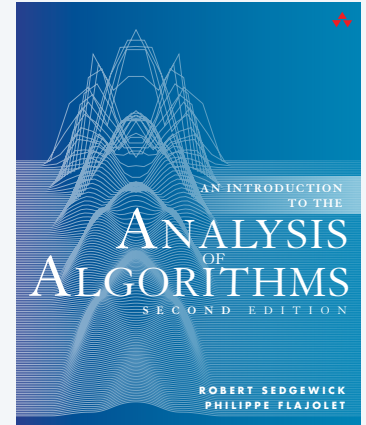


2. Run experiments to validate mathematical results.



Experiment 1. Write a program to generate and draw random tries (see lecture on Trees) and use it to draw 10 random tries with 100 nodes.

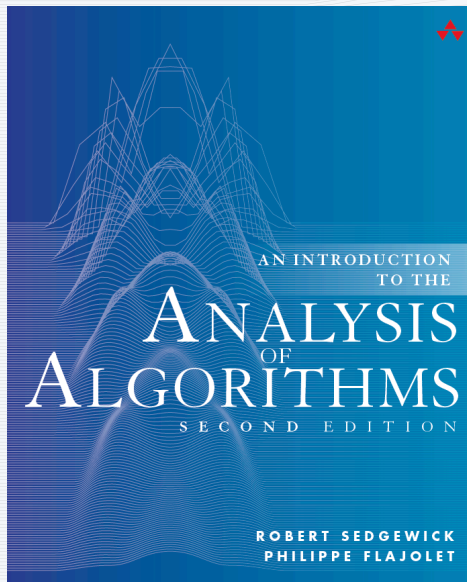
Experiment 2. *Extra credit.* Validate the results of the trie path length analysis by running experiments to build 100 random tries of size N for $N = 1000, 2000, 3000, \dots, 100,000$, producing a plot like Figure 1.1 in the text. Build the tries by inserting N random strings into an initially empty trie.



3. Write up solutions to Exercises 8.3, 8.14, and 8.57.

ANALYTIC COMBINATORICS

PART ONE



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8. Strings and Tries